## On the Convergence to a Mean-Field Equilibrium

#### Josu Doncel Polaris group INRIA

#### joint work with Nicolas Gast, Bruno Gaujal

Atelier en Évaluation de Performances Toulouse

16/03/2016

## Nash Equilibrium

The set of decisions such that no player has benefit of unilateral deviation

#### Theorem (Nash)

There exists a solution for every non-cooperative game.

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#### Theorem (Papadimitriou et al)

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#### $\Rightarrow$ Huge number of players or strategies: not tractable

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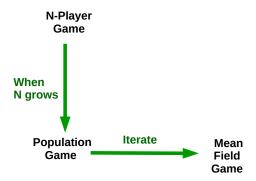
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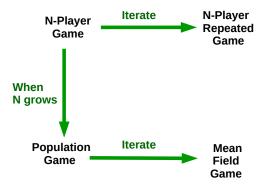
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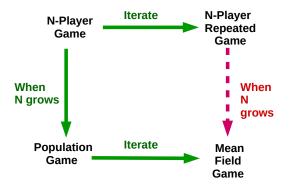
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 $\Rightarrow$  Huge number of players or strategies: not tractable

# Mean-Field Games (Lions and Lasry) Infinite number of rational objects in interaction. Josu Doncel (INBIA) MEG Convergence 16/03/2016









## 2 Convergence Results





## Mean-Field Games and Repeated Games

## 2 Convergence Results

#### 3 Some Extensions

## 4 Conclusions

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Infinite Homogeneous Players

- State:  $S = \{1, ..., S\}$
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Player 0

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$$a^0 \in \mathcal{A}$$

Instan. cost: c(a<sup>0</sup>(t), m<sup>a</sup>(t))

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$$V(a^0, a) = \mathbb{E}\left[(1 - \delta) \sum_{t=0}^{\infty} \delta^t \cdot c(a^0(t), \mathbf{m}^a(t))\right]$$

#### Best Response to a

Set of strategies

$$BR(a) = \operatorname*{arg\,min}_{a^0} V(a^0, a)$$

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#### Mean-Field Equilibrium

$$a^{\textit{MFE}} \in \textit{BR}(a^{\textit{MFE}})$$

 $\Rightarrow$  Fixed-point problem

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## N-Player Repeated Game

- Player  $0 \Rightarrow a^0$
- Others  $\Rightarrow a$

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#### N-player game equilibrium

For all a'

$$V^N(a,a) \leq V^N(a',a)$$

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#### Continuity assumptions on ${\boldsymbol{m}}$

- *P<sub>ija</sub>*(**m**)
- c(a<sup>0</sup>, **m**)

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#### Continuity assumptions on m

- *P<sub>ija</sub>*(**m**)
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#### Theorem

Any discrete time mean-field game with discounted cost that satisfies the continuity assumptions has a mean-field equilibrium.

Best-response has a fixed-point: Kakutani

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#### Player Strategies of N-player repeated game

- Local: depend on state and time  $\Rightarrow a(i, t)$
- Markov: depend on **m**, state and time  $\Rightarrow$  *a*(*i*, **m**(*t*))

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#### Theorem

Every local equilibrium converges to a mean-field equilibrium.

**Proof:**  $V^N(\pi',\pi) \to V(\pi',\pi)$  when  $N \to \infty$ 

(H. Tembine, J.-Y. L. Boudec, R. El-Azouzi, and E. Altman. Mean-field asymptotics of markov decision evolutionary games and teams. GameNets' 09.)

## Prisoner's dilemma

$$\mathcal{S} = \mathcal{A} = \{\mathcal{C}, \mathcal{D}\}$$

	С	D
С	1,1	3,0
D	0,3	2,2

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N Players: mean-field version  $\Rightarrow m_C + m_D = 1$ 

$$c(i,m) = \begin{cases} m_C + 3m_D & \text{if } i = C \\ 2m_D & \text{if } i = D \end{cases}$$

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Static game equilibrium  $\Rightarrow$  Always D

- MFE?
- Repeated Game NE?

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The best response to any action *a* is always D.

$$V(i,m) = (1-\delta) \sum_{t=0}^{\infty} \delta^{t} (x_{C}(t) \cdot (m_{C}(t) + 3m_{D}(t)) + x_{D}(t) \cdot 2m_{D}(t))$$

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Minimum when  $x_C(t) = 0 \Rightarrow BR(a)$ : Always D

## MFE: $a^D$ $a^D \in BR(a^D)$ Josu Doncel (INRIA) MFG Convergence 16/03/2016 13/21

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No player gets benefit from unilateral deviation

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- From k, if play D, an immediate advantage ⇒ punished until the end

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Let G be a symmetric matricial game, and let  $V^*$  be the cost under the strategy that repeats the Nash equilibrium of the static game G. Then any feasible cost V smaller than  $V^*$  is the cost of an equilibrium of the repeated game if the discount factor is large enough.

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### Why?

MFG assumption: Individuals action does not influence the mass

The mass can not punish

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#### The Folk Theorem do not scale

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MFG Convergence

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# 2 Convergence Results

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 $Q_{ija}(\mathbf{m})$ : transition rate matrix Assump:  $Q_{ija}(\mathbf{m})$  continuous in  $\mathbf{m}$ 

#### Theorem

Under the continuity assumptions, a MFE exists.

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Asynchronous N-player repeated game

- $T_N = \{0, 1/N, 2/N, ...\}$
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In continuous time, under the continuity assumptions and if  $\mathbb{E}(|R(t)|^2) < \infty$ , every local strategy converge to a MFE

Markov strategies

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MFF  $\Rightarrow a^D$ 

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In continuous time, under the continuity assumptions and if  $\mathbb{E}(|R(t)|^2) < \infty$ , every local strategy converge to a MFE

Markov strategies do not converge  $\Rightarrow$  Prisoner's dilemma

$$c(i, m^a) = \left\{ egin{array}{c} m_C + 3m_D & ext{if } i = 0 \ 2m_D & ext{if } i = L \end{array} 
ight.$$

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MFF  $\Rightarrow a^D$ 

 $\pi(m) = \begin{cases} C & \text{if } m_C = 1 \\ D & \text{if } m_C < 1 \end{cases}$ 

# Continuous Time and Finite Horizon

#### Theorem

Under the continuity assumptions, a MFE exists.

Previous example  $\Rightarrow$  MFE=NE (Always D)

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Ρ	0,3	0,4	3,3

Static Nash Equilibria: D and P

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Static Nash Equilibria: D and P

Repeated Game Nash Equilibrium:

- if t < 1, play C if  $m_C = 1$ , play P otherwise
- if  $t \ge 1$ , play *D* if  $m_P = 0$ , play *P* otherwise,

## 1 Mean-Field Games and Repeated Games

# 2 Convergence Results

## 3 Some Extensions



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- $\bullet \ \ \text{Local strategy} \rightarrow \text{MFE}$
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Discrete and continuous time

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MFE existence:

- Discrete and continuous time
- Discounted and finite horizon

• Best class of actions  $\Rightarrow$  convergence

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• Best class of actions  $\Rightarrow$  convergence

MFE:

Uniqueness conditions

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• Best class of actions  $\Rightarrow$  convergence

MFE:

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#### Full version available:

Josu Doncel, Nicolas Gast, Bruno Gaujal. Mean-Field Games with Explicit Interactions. https://hal.inria.fr/hal-01277098