

# Product form steady-state distribution for Stochastic Automata Networks

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## Stochastic Automata Networks in CT and DT

- To describe Multidimensional Markov Chains (Queueing Network, Stochastic Petri nets, SPA)
- The main idea is to obtain a generalized tensor description of the transition probability matrix or the transition rate matrix
- $N$  finite automata. One automaton is used to model one component.
- The state space is included into the Cartesian product of the state space of the automata.
- The links of the automata carry information:
  - rate (in CT) or probability (DT) : fixed or function
  - local or synchronization

## SAN and Tensor

- For CT-SAN, transition rate matrix is given by:

$$Q = \bigoplus_g^N Q_l^i + \sum_s \bigotimes_g^N Q_s^{(i)} + D$$

where  $D$  is a diagonal matrix (for normalization),  $\bigoplus_g$  and  $\bigotimes_g$  are the generalized tensor sum and the generalized tensor product and  $Q_l^{(i)}$  and  $Q_s^{(i)}$  are matrices describing the local transitions and transitions due to synchronization  $s$  on automaton  $i$ .

- Also proved for many formalisms (PEPA, Petri nets).
- for models without functions, use  $\otimes$  and  $\oplus$  instead of  $\otimes_g$  and  $\oplus_g$ .
- For DT-SAN, transition probability matrix is given by:

$$P = \bigotimes_g^N P_l^i + \sum_s \bigotimes_g^N P_s^{(i)} + D$$

## Generalized Tensor Product and Sum

- Ordinary tensor product  $C = A \otimes B$  is defined by assigning the element of  $C$  that is in the  $(i_2, j_2)$  position of block  $(i_1, j_1)$ , the value  $a_{i_1 j_1} b_{i_2 j_2}$ . We shall write this as

$$c_{\{(i_1, j_1); (i_2, j_2)\}} = a_{i_1 j_1} b_{i_2 j_2}.$$

- Generalized Tensor Product: Matrices of functions whose arguments are the states of the other components (ie the index of the matrix).

$$c_{\{(i_1, j_1); (i_2, j_2)\}} = a_{i_1 j_1}(i_2) b_{i_2 j_2}(i_1),$$

- As usual the sum is defined using the product:

$$D = A(\mathcal{B}) \oplus_g B(\mathcal{A}) \Leftrightarrow D = A(\mathcal{B}) \otimes_g Id_B + Id_A \otimes_g B(\mathcal{A}),$$

## Compatibility between $\otimes$ (or $\otimes_g$ ) and $\times$

- associativity and distributivity: OK for  $\otimes$  and  $\otimes_g$ .
- Most important property: Compatibility between ordinary product and ordinary tensor product.

$$(A \otimes B) \times (C \otimes D) = (A \times C) \otimes (B \times D)$$

- Imagine that  $A$  is not a matrix but a vector  $v \dots$
- What if  $v$  is in the kernel (null space) of  $C \dots$  or  $v$  is an eigenvector of  $C$  ?
- But the property is not valid for  $\otimes_g$  in general.
- $0 \otimes v = 0$ .

## Product Form

- The steady-state distribution is the product of marginal distributions.

$$\pi(x_1, x_2, \dots, x_n) = C \pi_{x_1} \pi_{x_2} \dots \pi_{x_n}$$

- Proved for some queueing networks (Jackson, BCMP, G-networks), Petri-Nets, SAN, PEPA models
- First remark : a product of marginal distributions is a tensor product of vectors.

## Motivation

- Links between (generalized or ordinary) tensor decomposition of multidimensional CTMC and product form for steady-state distribution.
- Is it possible for  $M = \sum_i \otimes_{g_j} P_j^i$  to have a steady-state solution  $\pi = \otimes_j \pi_j$  where  $\pi_j$  are small vectors associated with  $P_j^i$  ?

- 

$$(\otimes_j \pi_j) \times \left( \sum_i \otimes_{g_j} P_j^i \right) = 0$$

- Infinite State Space.

## A little bit of history

- CT-Models with functions but without synchronization  
$$Q = \bigoplus_{i=1}^N Q^i$$

Generalization of many results (Hillston, Boucherie, Robertazzi) for Stochastic Petri Nets, Markov Chain in Competition, Modulated queues (ValueTools07, Performance Evaluation)
- Much harder in Discrete Time but still possible (Qest08)
- With synchronizations but without functions. First, give the number of components involved in an arbitrary synchronization:
  - 2: Master/Slave (like in a Jackson queueing network (PAPM 97), here with a new proof)
  - 3: Domino, like a G-network with triggered customer movements (Epew 2008).
- Properties of the generalized tensor product (Model 35)



## Part I : Continuous-Time

### Master/slave synchronizations, no functions

$$Q = \bigoplus_{i=1}^N Q_l^i + \sum_s \bigotimes_{i=1}^N Q_s^{(i)} + D$$

## Master-Slave: Matrix Description

- The master initiates the synchronization.
- The slave follows.
- A synchronization is described by 2 matrices.
  - Master description:  $M^{(r)}$ , a transition rate matrix.
  - Slave description:  $E^{(r)}$ , a transition probability matrix
- Local Transitions: transition rate matrices  $F_l$ .

## Normalization

- All matrices are normalized, i.e. for all  $k$  we have:

$$M^{(r)}[k, i] \geq 0 \text{ if } i \neq k \text{ and } \sum_i M^{(r)}[k, i] = 0,$$

$$E^{(r)}[k, i] \geq 0 \text{ and } \sum_i E^{(r)}[k, i] = 1,$$

- Normalization of the SAN: based on the normalization of the synchronizations ( $N^r$ ) as the local transitions are already associated to transition rate matrices.
- Definition: Let  $M$  be a matrix,  $diag(M)$  is a diagonal matrix whose elements are the diagonal elements of  $M$ .

## A normalized tensor representation of a Master-Slave

1.  $(M^{(r)} - \text{diag}(M^{(r)})) \otimes E^{(r)} \otimes I_1$ : the slave accepts the synchronization.
2.  $\text{diag}(M^{(r)}) \otimes I \otimes I_1$ : normalization of term 1.

## Main Result

Theorem: Consider a SAN with  $n$  automata and  $s$  with Master/Slave synchronizations. Consider matrices  $\overline{M}^{(r)} = M^{(r)} - \text{diag}(M^{(r)})$  and  $E^{(r)}$  associated to the description of synchronization  $r$ . Let  $g_l$  an eigenvector of  $\overline{M}^{(r)}$ . We assume that  $g_l > 0$ . Let  $\Gamma_r$  be the eigenvalue for matrix  $\overline{M}^{(r)}$  associated to  $g_l$ .

If  $g_l$  is in the kernel of matrix

$$A_l = F_l + \sum_{r=1}^R \left( M^{(r)} 1_{msr(r)=l} + \Gamma_r (E^{(r)} - I) 1_{sl(r)=l} \right),$$
 then the

steady-state distribution has a product form solution:

$$Pr(X_1, X_2, \dots, X_n) = C \prod_{l=1}^n g_l(X_l), \quad (1)$$

and  $C$  is a normalization constant.

## Proof by tensor algebra

$(\otimes_l g_l) \times (\oplus_l F_l + \sum_i \otimes_l B_l^i)$	Reorganize the product (distributivity)
$\sum_i (\otimes_l g_l) \times (\otimes_l M_l^i)$	Then use compatibility with $\times$
$\sum_i (\otimes_l (g_l \times M_l^i))$	Assumptions on eigenvectors
	Factorize and find some matrices $P_l^i$
$\sum_i (\otimes_l (g_l \times P_l^i))$	such that
$\pi_{j_0} P_{j_0}^i = 0$	for all $i$ , it exists $j_0$ (automaton index)

## Proof-2

The description of  $(\pi_1 \otimes \pi_2 \otimes \dots \otimes \pi_n)Q$  consists in 4 terms (two coming from the tensor sum, one for the Master/Slave and one for the normalization of the Master/Slave):

$$\begin{aligned} & (g_1 F_1 \otimes g_2 \otimes \dots \otimes g_n) \\ + & (g_1 \otimes g_2 F_2 \otimes \dots \otimes g_n) \\ + & (g_1 (M^{(r)} - \text{diag}(M^{(r)})) \otimes g_2 E^{(r)} \otimes \dots \otimes g_n) \\ + & (g_1 \text{diag}(M^{(r)}) \otimes g_2 I \otimes \dots \otimes g_n) \end{aligned}$$

## Proof-3

- Now remember that  $g_1(M^{(r)} - \text{diag}(M^{(r)})) = g_1\Gamma_r$ . And of course  $g_2I = g_2$ . After simplification, we get:

$$\begin{aligned} & (g_1F_1 \otimes g_2 \otimes \dots \otimes g_n) \\ + & (g_1 \otimes g_2F_2 \otimes \dots \otimes g_n) \\ + & (g_1\Gamma_r \otimes g_2E^{(r)} \otimes \dots \otimes g_n) \\ + & (g_1\text{diag}(M^{(r)}) \otimes g_2 \otimes \dots \otimes g_n) \end{aligned}$$



## Proof-4

- Now, remark that  $g_1\Gamma_r \otimes g_2E^{(r)} = g_1 \otimes g_2\Gamma_r E^{(r)}$  because the ordinary product is compatible with the tensor product.
- We factorize the first and the last terms and we do the same for the second and the third term. Furthermore we add and subtract the following term:  $(g_1(M^{(r)} - \text{diag}(M^{(r)})) \otimes g_2 \otimes \dots \otimes g_n)$ .

$$\begin{aligned} & (g_1(F_1 + \text{diag}(M^{(r)})) \otimes g_2 \otimes \dots \otimes g_n) \\ + & (g_1 \otimes g_2(F_2 + \Gamma_r E^{(r)}) \otimes \dots \otimes g_n) \\ - & (g_1(M^{(r)} - \text{diag}(M^{(r)})) \otimes g_2 \otimes \dots \otimes g_n) \\ + & (g_1(M^{(r)} - \text{diag}(M^{(r)})) \otimes g_2 \otimes \dots \otimes g_n) \end{aligned}$$

## Proof-5

- We factorize the first and the last term and we note that  $g_1(M^{(r)} - \text{diag}(M^{(r)})) = g_1\Gamma_r$  to simplify the third term:

$$\begin{aligned} & (g_1(F_1 + M^{(r)}) \otimes g_2 \otimes g_3 \otimes \dots \otimes g_n) \\ + & (g_1 \otimes g_2(F_2 + \Gamma_r E^{(r)}) \otimes \dots \otimes g_n) \\ - & (g_1\Gamma_r \otimes g_2 \otimes \dots \otimes g_n) \end{aligned}$$

- Again we use the compatibility of the ordinary product with the tensor product and we get after factorization:

$$\begin{aligned} & (g_1(F_1 + M^{(r)}) \otimes g_2 \otimes g_3 \otimes \dots \otimes g_n) \\ + & (g_1 \otimes g_2(F_2 + \Gamma_r(E^{(r)} - I)) \otimes \dots \otimes g_n) \end{aligned}$$

- This is the decomposition we need.

## Example-1: G-network with + and - customers

- Infinite state space.
- Each automaton models the number of positive customers in a queue.
- The local transitions are the external arrivals (rate  $\lambda_l$ ) and the departures to the outside (rate  $\mu_l$  multiplied by probability  $d_l$ ).
- Synchronization : departure of a customer on the master (the end of service with rate  $\mu_l$  and probability  $(1 - d_l)$ ), the departure of a customer on the slave if there is any.

## Example-1-Matrix

$$\begin{aligned}Q^{(l)} &= \lambda_l(U - I) + \mu_l d_l(L - I^0), \\M^{(r)} &= \mu_l(1 - d_l)(L - I^0), \\E^{(r)} &= L\end{aligned}\tag{2}$$

- Tridiagonal matrix.
- $\pi_l$  has a geometric distribution with rate  $\rho_l$ :

$$\rho_l = \frac{\lambda_l}{\mu_l + \sum_{r=1}^R \Gamma_r 1_{sl(r)=l}}.$$

- $\pi_l$  is an eigenvector of operators  $\overline{M^{(r)}}$  and  $E^{(r)}$ .
- Finally:  $\Omega_r = \rho_{sl(r)}$  and  $\Gamma_r = \rho_{msr(r)} \mu_{msr(r)} (1 - d_{msr(r)})$ ,

## Part II : Continuous-Time

functional rates, no synchronizations

$$Q = \bigoplus_{i=1}^N Q_l^i$$

## Here

- Infinite State Space.
- No Synchronizations.
- Functions to model the interactions between components.
- An easy model to represent multidimensional Markov chains:
  - without synchronized transition: only one component change during a transition.
  - with transition rates which are functions of the other components.
- $Q^{(l)}[k_l, i](\vec{k}, \vec{k} + (l, i))$  : transition rate matrix for automaton  $l$ . The state of the automaton jumps from  $k_l$  to  $i$ . Due to this local jump, the global state changes from  $\vec{k}$  to  $\vec{k} + (l, i)$ . The rate may depend of the global state (i.e. fonctionnal rate).

## Example

- Consider a SAN with two automata  $\mathcal{A}_1$  and  $\mathcal{A}_2$ .
- Both have a very simple state space:  $\{0, 1\}$
- The transitions in  $\mathcal{A}_1$  have a fixed rate  $l_1$  for the transition from 0 to 1 and  $l_2$  for the transition from 1 to 0.
- Automaton  $\mathcal{A}_2$  has two functional transitions: the rate from 0 to 1 has a functional rate  $f_0$  and the reverse transitions has functional rate  $f_1$ . Both functions use the state of automaton  $\mathcal{A}_1$  as an argument (denoted as  $x_1$ ).
- $f_0(x_1) = mb + m(1 - b)1_{x_1=0}$  and  $f_1(x_1) = m_1 + m_2 1_{x_1=0}$ .

## Example-SAN

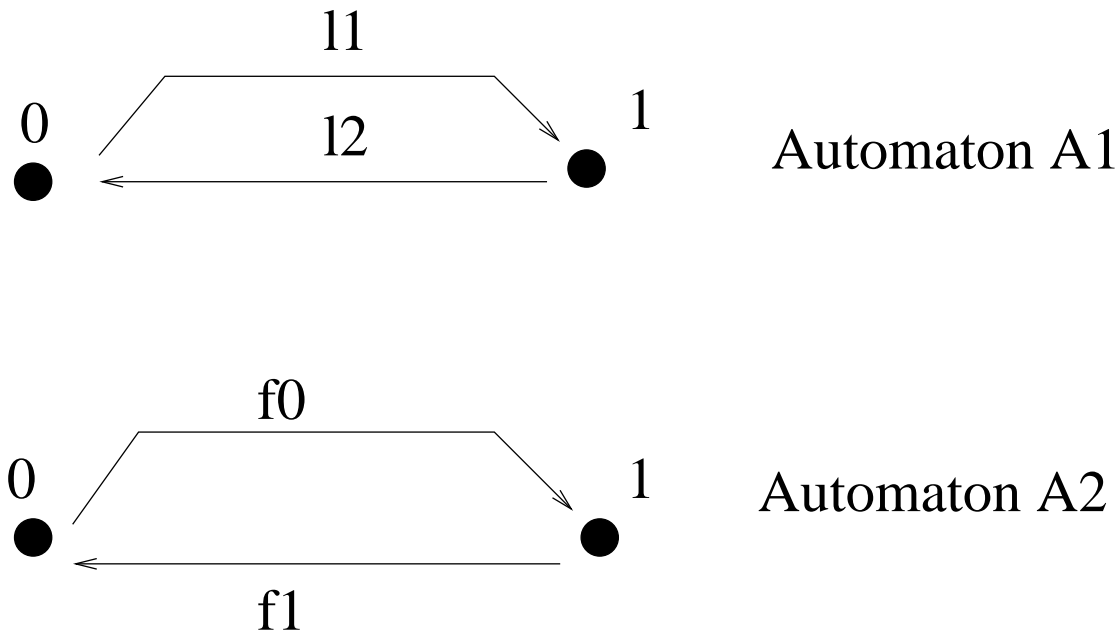


Figure 1: Stochastic Automata Network



## Example

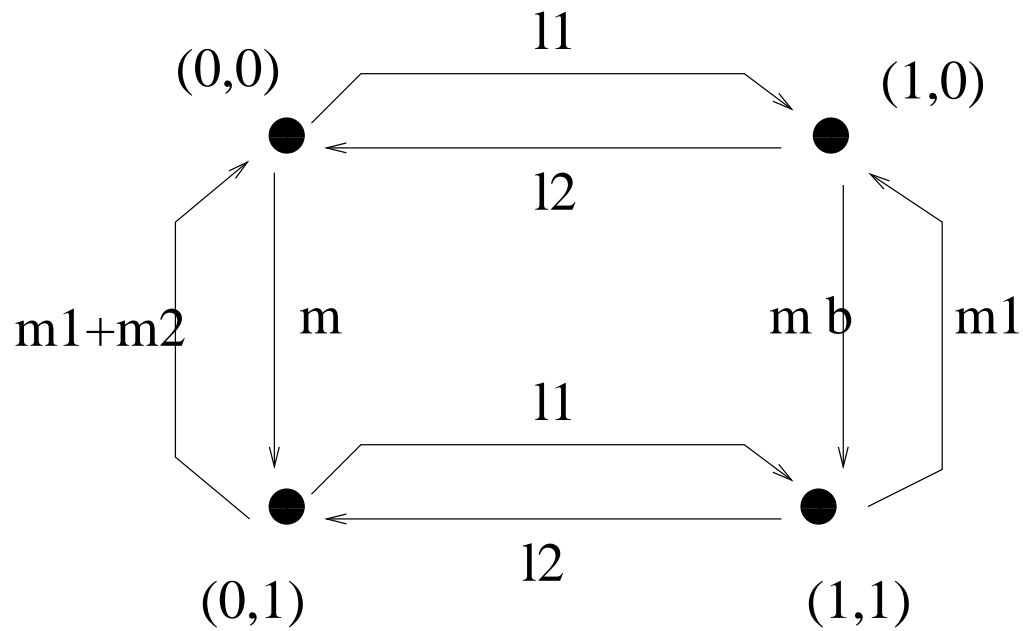


Figure 2: Markov chain

## CK equation

$$Pr(\vec{k}) \left[ \sum_{l=1}^n \sum_{i \neq k_l} Q^{(l)}[k_l, i](\vec{k}, \vec{k} + (l, i)) \right] =$$

(3)

$$\sum_{l=1}^n \sum_{i \neq k_l} Q^{(l)}[i, k_l](\vec{k} + (l, i), \vec{k}) Pr(\vec{k} + (l, i)).$$

## Main Idea

- As the state space is discrete, functions can be replaced by an index.
- **Definition 1** *Let  $l$  be an automaton index, we consider all the functions in matrix  $Q^{(l)}$  and we evaluate them for all state  $\vec{k}$  when the transition from  $\vec{k}$  to  $\vec{k} + (l, i)$  takes place. Such a matrix will be denoted by  $L^{(l, m(\vec{k}))}$  where  $m(\vec{k})$  is an index. The set of matrices  $L^{(l, m(\vec{k}))}$  will be denoted by  $\mathcal{F}_{(l)}$ .*

## Definition


- **Definition 2** *Let  $\alpha$  be a probability distribution. We note by  $\mathcal{S}(\alpha)$  the set of transition rate matrices  $M$  such that  $\alpha M = 0$  (i.e.  $\alpha$  is in the kernel of all matrices in  $\mathcal{S}(\alpha)$ ).*
- **Property 1** *Interesting properties of  $\mathcal{S}(\alpha)$ :*
  1.  $\mathbf{0}$  (the matrix whose elements are all zero) is in  $\mathcal{S}(\alpha)$
  2.  $aM1$  is in  $\mathcal{S}(\alpha)$ . for all matrices  $M1$  in  $\mathcal{S}(\alpha)$  and  $a$  in  $R^+$ .
  3.  $aM1 + bM2$  is in  $\mathcal{S}(\alpha)$  for all matrices  $M1$  and  $M2$  in  $\mathcal{S}(\alpha)$  and  $a, b$  in  $R^+$  such that  $a + b = 1$ .

## Main Theorem

- **Theorem 1** *Consider a SAN with functions but without synchronizations. Assume that the steady state exists. If for each automaton  $l$  there exists a probability distribution  $\pi_l$  such that all the matrices in  $\mathcal{F}_{(l)}$  are in  $\mathcal{S}(\pi_l)$ , then the SAN has a product form steady state distribution such that:*

$$Pr(x_0, \dots, x_n) = C \pi_1(x_1) \dots \pi_l(x_l) \pi_n(x_n).$$

- The proof is based on the resolution of the Chapman-Kolmogorov equation at steady-state.



$$Pr(\vec{k}) \left[ \sum_{l=1}^n \sum_{i \neq k_l} L^{(l, m(\vec{k}))} [k_l, i] \right] =$$

$$\sum_{l=1}^n \sum_{i \neq k_l} L^{(l, m(\vec{k}))} [i, k_l] Pr(\vec{k} + (l, i)) \quad . \quad (4)$$

**Corollary 1** *Consider the previous example. Matrices  $M_0$  and  $M_1$  have the same kernel if  $b = \frac{m_1}{m_1 + m_2}$ . If this condition is satisfied, the steady-state distribution of the SAN has product form:*

$$\pi(x_1, x_2) = C \left( \frac{l_1}{l_2} \right)^{x_1} \left( \frac{m}{m_1 + m_2} \right)^{x_2} .$$

## Irreducibility

- The CTMC must be irreducible.
- It is sufficient that the components are irreducible but it is not necessary.
- If some matrices in  $\mathcal{S}(\alpha)$  are reducible we can still obtain an irreducible CTMC with product form.

## Example with reducible matrices

- A network with two automata  $\mathcal{A}_1$  and  $\mathcal{A}_2$ .
- $\mathcal{A}_2$  has a very simple state space:  $\{0, 1\}$
- $\mathcal{A}_1$  has four states
- The transitions in  $\mathcal{A}_2$  have a fixed rate  $l_2$  for the transition from 0 to 1 and  $m_2$  for the transition from 1 to 0. Automaton  $\mathcal{A}_1$  contains four functional transitions governed by four functions  $f_1$ ,  $f_2$ ,  $f_3$  and  $f_4$ .



## Example - SAN

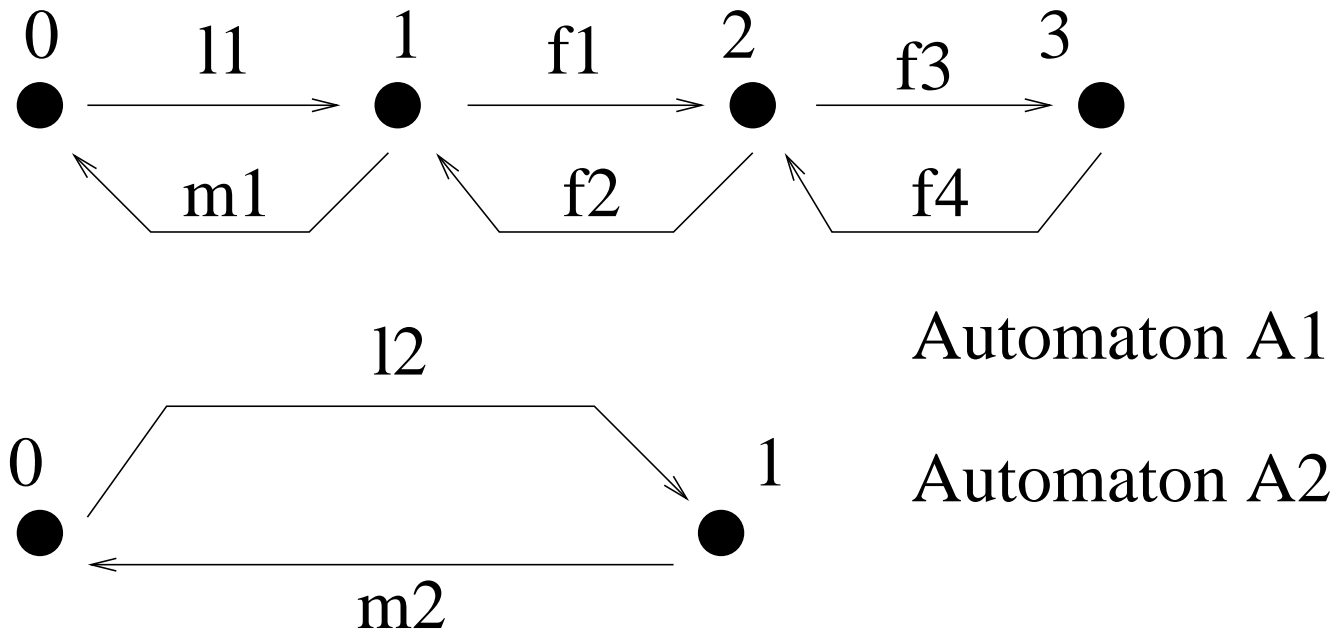


Figure 3: A more complex SAN with product form.

## Example - Matrices

- $$\begin{cases} f1 = l1 & 1_{x2=1} & f2 = m1 & 1_{x2=1} \\ f3 = l1 & 1_{x2=0} & f4 = m1 & 1_{x2=0} \end{cases}$$

- $$M0 = \begin{pmatrix} -l1 & l1 & 0 & 0 \\ m1 & -m1 & 0 & 0 \\ 0 & 0 & -l1 & l1 \\ 0 & 0 & m1 & -m1 \end{pmatrix} \quad \text{and} \quad M1 = \begin{pmatrix} -l1 & l1 & 0 & 0 \\ m1 & -m1 - l1 & l1 & 0 \\ 0 & m1 & -m1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} .$$

## Example - Kernels

- $M0 : \{u (m1, l1, 0, 0) + v (0, 0, m1, l1), \forall u, v \in \mathcal{R}\}$
- $M1 : \{u (1, l1/m1, l1^2/m1^2, 0) + v (0, 0, 0, 1), \forall u, v \in \mathcal{R}\}.$
- Clearly the vector  $(1, l1/m1, l1^2/m1^2, l1^3/m1^3)$  is in both sets.
- Product form solution.

## Tensor based proof

- The Hidden Lemma :
- **Property 2** *Let  $A(\mathcal{B})$  and  $B(\mathcal{A})$  be arbitrary functional transition rate matrices. Assume that  $w$  is in the kernel of  $B(y)$  for every  $y$  and that  $w$  is positive. Similarly assume that there exists a positive vector  $v$  which is in the kernel of  $A(x)$  for all  $x$ . Then we have:*

$$(v \otimes w) \times (A(\mathcal{B}) \oplus_g B(\mathcal{A})) = 0.$$

- Very simple proof (algebra).

## Previous results

- Plateau's first theorem on product form for SAN
- Boucherie's first theorem on competing Markov chains.
- Verchere's theorem on modulated Markov Chains
- Partial Reversibility
- They are all corollaries of our main theorem.

## Plateau's first theorem on Product Form SAN

- SAN with functions.
- The transition rate matrix of automaton  $l$  is the product of a function of  $\vec{k}$  except component  $l$  ( $f_l(\vec{k})$ ) by an usual transition rate matrix.
- $Q^{(l)}[i, k_l](\vec{k} + (l, i), \vec{k}) = f_l(\vec{k})Q^{(l)}[i, k_l]$
- All these matrices have the same dominant eigenvector.

## Boucherie's first theorem on competing MC

- Associated to Petri nets.
- A collection of Markov chains and a product process with restriction on the state space.
- Competition over ressources.
- Uniformally: if a ressource is owned by component (i.e. a chain), transitions from some other chains (i.e. the competing ones) are removed.

## Example

- Two chains  $X1$  and  $X2$  both with states  $\{0, 1, 2, 3\}$  competing over one resource.
- Symmetrical rules.
- The resource is owned by a chain when it is in state 2 or 3.
- It is released when the chain jumps from state 3 to 1.
- Thus states in  $\{2, 3\} \times \{2, 3\}$  are forbidden.
- When process  $X1$  is in state 2 or 3 process  $X2$  is stopped. If process  $X1$  is in state 0 or 1, process  $X2$  can move.



# Graph of the example

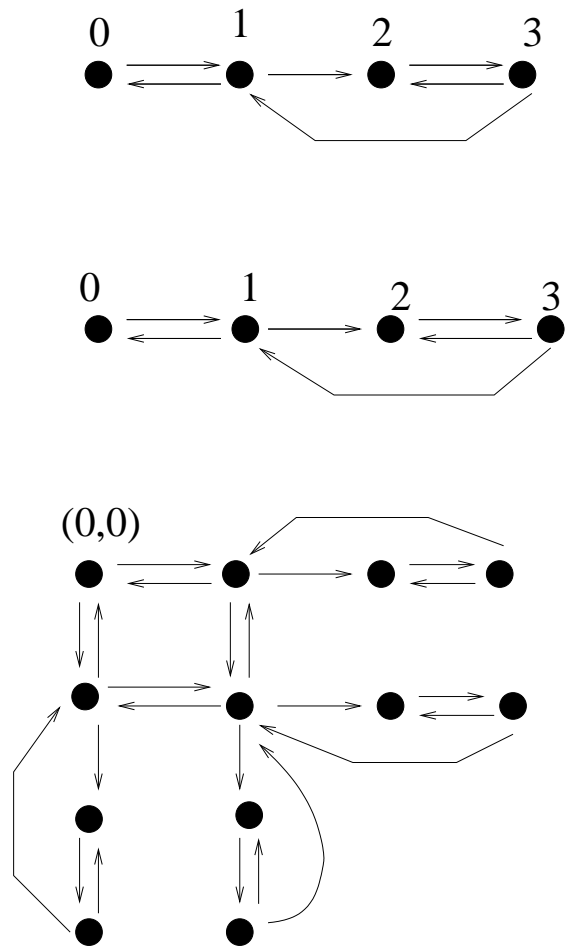


Figure 4: Two Markov chains in competition

## Transitions of a competing Markov chain

- if states  $\vec{k}$  and  $\vec{k}'$  differ by more than 1 components, the transition rate is 0. (the transition matrix is a tensor sum of some matrices).
- from state  $\vec{k}$  to state  $\vec{k} + (l, i)$  the transition is the transition rate from  $k_l$  to  $i$  in chain  $l$  multiplied by an indicator function.
- This function is equal to zero when there exists a resource  $r$  owned by another chain which competes with  $l$ . (the transition rate matrices are the original matrices of the chains multiplied by a function of the states which takes value in  $\{0, 1\}$ ).
- **This is a simple corollary of Plateau's first theorem where the functions take value in  $\{0, 1\}$ .**

## Partial Reversibility

- **Definition 3** *An ergodic Markov chain  $W$  (matrix  $F$ ) is partially reversible if and only if there exists a non empty subset  $X$  of the states such that for all states  $i$  and  $j$  in  $X$  we have local balance equations between  $i$  and  $j$ :  $\pi(i)F(i, j) = \pi(j)F(j, i)$ .*
- **Property 3** *Assume  $F \in \mathcal{S}(\alpha)$ ,  $F$  irreducible and partially reversible with set  $X$ , then all matrices obtained from  $F$  after*
  1. *choosing any subset of  $X$  (let call it  $Y$ )*
  2. *multiplying all the transitions between states of  $Y$  by an arbitrary positive constant  $c$ .*

*are also in  $\mathcal{S}(\alpha)$ .*

*We delete the transitions but we do not delete the states.*

## Partial Reversibility and Product Form

**Theorem 2** *Consider a SAN with two automata  $\mathcal{A}_1$  and  $\mathcal{A}_2$ . Assume that:*

- 1. all the rates are constant in  $\mathcal{A}_1$ ,*
- 2. the matrix of  $\mathcal{A}_2$  is partially reversible, ( $X$ : set of states of  $\mathcal{A}_2$  with local balance,  $Y$ : an arbitrary subset of  $X$ ,  $R$ : matrix of the transitions between states of  $Y$ ),*
- 3.  $\mathcal{A}_2$  contains functions whose argument is the state of  $\mathcal{A}_1$ , and the functions are only carried by the transition between the states of  $Y$ .*
- 4.  $R$  is the product of a function  $f$  by a constant matrix  $R_0$ ,*
- 5. the CTMC is ergodic,*

*then the SAN has a product form solution.*

## Checking Partial Reversibility

- Sometimes due to structural properties
- **Definition 4 (peninsula)** *Consider an ergodic CTMC associated to transition rate matrix  $F$ , a peninsula is a set of two nodes  $a, b$  such that:*
  - *Removing  $a$  and  $b$  disconnects the chain and creates two connected components  $A$  and  $B$ .*
  - *$a \in A$  and  $b \in B$ .*
  - *$b$  is the only one successor of  $a$ .*
  - *$a$  is the only one successor of  $b$ .*
- A peninsula implies a local balance between  $a$  and  $b$ .

## Modulated network of queues

- One automata to represent the phase and one to represent the network of queues.
- Thus the synchronized transition between queues are local to the second automata.
- The transitions of the queues (not only the rate) may depend of the state of the phase.
- Verchère's theorem: if the steady-state distribution of the queueing network is always the same for all state of the phase, then the global system has a product form steady-state distribution.

## Not that simple

- A two state phase.
- In phase 1, we have a Jackson network (transition  $(-1,+1)$ ).
- In phase 2, a G-network with positive customers (transition  $(-1,+1)$ ), triggers (transition  $(-1,-1,+1)$ ) and negative customers (transition  $(-1,-1)$ ).
- Both networks do not have the same transitions (because of negative customers and triggers).
- But if the rates are carefully chosen, they have the same geometric steady-state distribution
- Product-form.

## Part III : Discrete-Time

fonctional rates, no synchronizations

$$P = \bigotimes_{i=1}^N P_l^i$$



## Simple models for DT

- Infinite State Space.
- Local Events.
- Functions to model the interactions between components.
- Several components change during a transition (DT).
- With transition probabilities which are functions of the other components.
- Discrete-Time:  $P = \bigotimes_{i=1}^N P_l^i$   
and  $P_l^i$  is a functional transition matrix
- remark that in Discrete-Time without function:  $P = \bigotimes_{i=1}^N P_l^i$  has a product-form steady-state solution (independence)...

## Functional Dependency Graph

- Functional Dependency Graph (FDG): directed graph  $(V, E)$
- Node = Automaton.
- Directed edges  $(A1, A2)$ : automaton  $A1$  uses the state of  $A2$  in some functions to define rates or probabilities.
- The numerical algorithm developed by Plateau, Stewart, and Fernandes takes into account some properties of the Functional Dependency Graph.

## Main Idea

- As the state space is discrete, functions can be replaced by an index.
- **Definition 5** *Let  $l$  be an automaton index, we consider all the functions in matrix  $Q^{(l)}$  and we evaluate them for all state  $\vec{k}$  when the transition from  $\vec{k}$  to  $\vec{k} + (l, i)$  takes place. Such a matrix will be denoted by  $L^{(l, m(\vec{k}))}$  where  $m(\vec{k})$  is an index. The set of matrices  $L^{(l, m(\vec{k}))}$  will be denoted by  $\mathcal{F}_{(l)}$ .*

## Definition

- **Definition 6** Let  $\alpha$  be a probability distribution. We note by  $\mathcal{S}(\alpha)$  the set of transition rate matrices  $M$  such that  $\alpha M = 0$  (i.e.  $\alpha$  is in the kernel of all matrices in  $\mathcal{S}(\alpha)$ ).
- The definition comes from CT models and it kept for compatibility reasons but it implies that we must transform transition probability matrix to use it.
- DT models: if two stochastic matrix  $M1$  and  $M2$  have the same dominant eigenvector  $\alpha$ , then both  $(M1 - Id)$  and  $(M2 - Id)$  are in  $\mathcal{S}(\alpha)$ .
- **Property 4** Interesting properties of  $\mathcal{S}(\alpha)$ :
  1.  $\mathbf{0}$  (the matrix whose elements are all zero) is in  $\mathcal{S}(\alpha)$
  2.  $aM1 + bM2$  is in  $\mathcal{S}(\alpha)$  for all matrices  $M1$  and  $M2$  in  $\mathcal{S}(\alpha)$  and  $a, b$  in  $R^+$  such that  $a + b = 1$ .

## Example

- Consider the functional matrix:

$$A(x) = \begin{bmatrix} 1 - x/2 & x/4 & x/4 \\ 0 & 1 - x/2 & x/2 \\ x & 0 & 1 - x \end{bmatrix}.$$

- If  $0 < x < 1$ ,  $A(x)$  is finite and irreducible.
- Steady-state distribution  $(1/2, 1/4, 1/4)$  and it does not depend of the value of  $x$ .

## Relations with the generalized tensor product

- **Property 5** *Let  $B$  be a positive matrix, let  $A(\mathcal{B})$  be a matrix whose elements are functions of the index of  $B$ . Assume that  $w$  is an eigenvector of  $B$  with eigenvalue  $\lambda$ . Assume that for all states  $s$  of  $B$ ,  $A(s)$  has an eigenvector  $v$  associated to eigenvalue  $\mu$ . Assume that both  $\mu$  and  $v$  do not depend of  $s$ . Then we have:*

$$(v \otimes w) \times (A(\mathcal{B}) \otimes_g B) = \lambda\mu (v \otimes w).$$

- Proof: simple algebra.
- Easy Generalization to an arbitrary number of automata....
- But Functional Dependency Graph = DAG...

## Main result for Discrete Time

**Theorem 3** *Consider a collection of functional stochastic matrices such that:*

- *The functional dependency graph is a Directed Acyclic Graph.*
- *For every matrix  $l$  there exists a positive vector  $\pi_l$  such that for every matrix index  $m$ ,  $\pi_l$  is in the kernel of matrix  $(Q^{(l,m)} - Id)$*
- *The Markov chain associated to the composition of these functional matrices is ergodic.*

*Then the steady-state distribution has product form.*

$$Pr(x_0, \dots, x_n) = C \pi_1(x_1) \dots \pi_l(x_l) \pi_n(x_n).$$

## What if the FDG is not a DAG

- Numerical Example.
- SAN with automata  $\mathcal{A}_1$  and  $\mathcal{A}_2$ .
- States of  $\mathcal{A}_1$  are labelled 1, 2 and 3.
- States of  $\mathcal{A}_2$  are labelled 1 and 2.
- Transition matrix of  $\mathcal{A}_1$  is functional matrix  $A(x)$  where  $x$  is the state of  $\mathcal{A}_2$  divided by 4.
- Transition matrix of  $\mathcal{A}_2$  is functional matrix  $D(y)$  where  $y$  is the state of  $\mathcal{A}_1$  divided by 2.
- The dependency graph is a complete directed graph on two states and this is not a DAG.



## Matrix $D(y)$

- $D(y) = \begin{bmatrix} 1 - y/2 & y/2 \\ y/4 & 1 - y/4 \end{bmatrix}$ .
- If  $y$  is between 0 and 2,  $D(y)$  is finite and irreducible.
- It has a steady-state distribution:  $(1/3, 2/3)$  and it does not depend of the value of  $y$ .

## Matrix of the SAN

- 

$$\begin{bmatrix} 0.65625 & 0.21875 & 0.046875 & 0.015625 & 0.046875 & 0.015625 \\ 0.09375 & 0.65625 & 0.015625 & 0.109375 & 0.015625 & 0.109375 \\ 0 & 0 & 0.4375 & 0.4375 & 0.0625 & 0.0625 \\ 0 & 0 & 0.1875 & 0.5625 & 0.0625 & 0.1875 \\ 0.0625 & 0.1875 & 0 & 0 & 0.1875 & 0.5625 \\ 0.1875 & 0.3125 & 0 & 0 & 0.1875 & 0.3125 \end{bmatrix}$$

- its steady-state distribution is

$$\pi = \left[ 0.1923 \quad 0.3169 \quad 0.0803 \quad 0.1664 \quad 0.0751 \quad 0.1690 \right],$$

- but the tensor product of the two steady-state distribution is:

$$\pi_A \otimes \pi_D = \left[ 1/6 \quad 1/3 \quad 1/12 \quad 1/6 \quad 1/12 \quad 1/6 \right]$$

- The SAN does not have a product form distribution

## What about $\oplus_g$

- **Property 6** *Let  $A(\mathcal{B})$  and  $B(\mathcal{A})$  be arbitrary functional transition rate matrices. Assume that  $w$  is in the kernel of  $B(y)$  for every  $y$  and that  $w$  is positive. Similarly assume that there exists a positive vector  $v$  which is in the kernel of  $A(x)$  for all  $x$ . Then we have:*

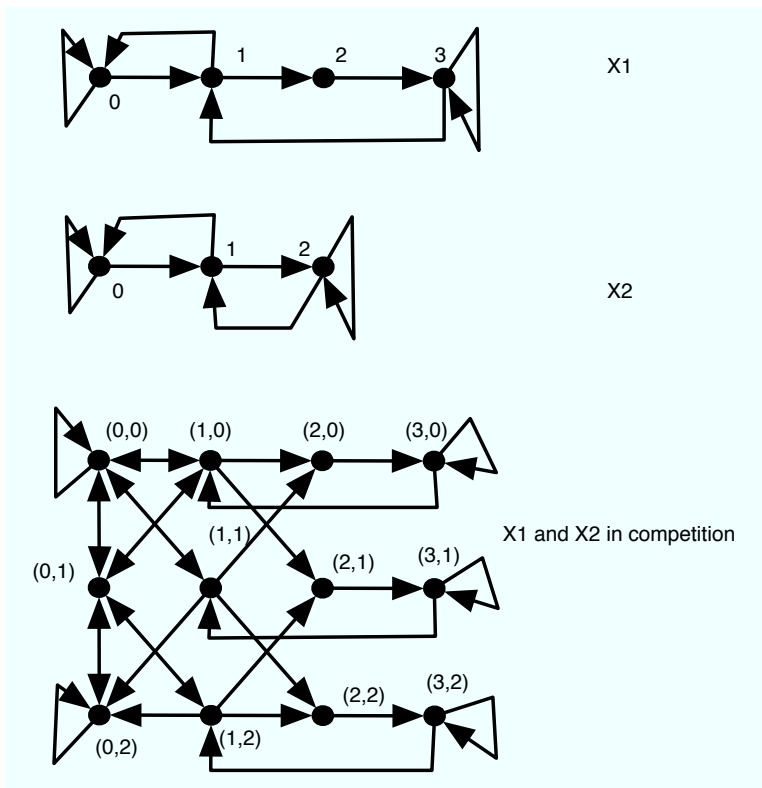
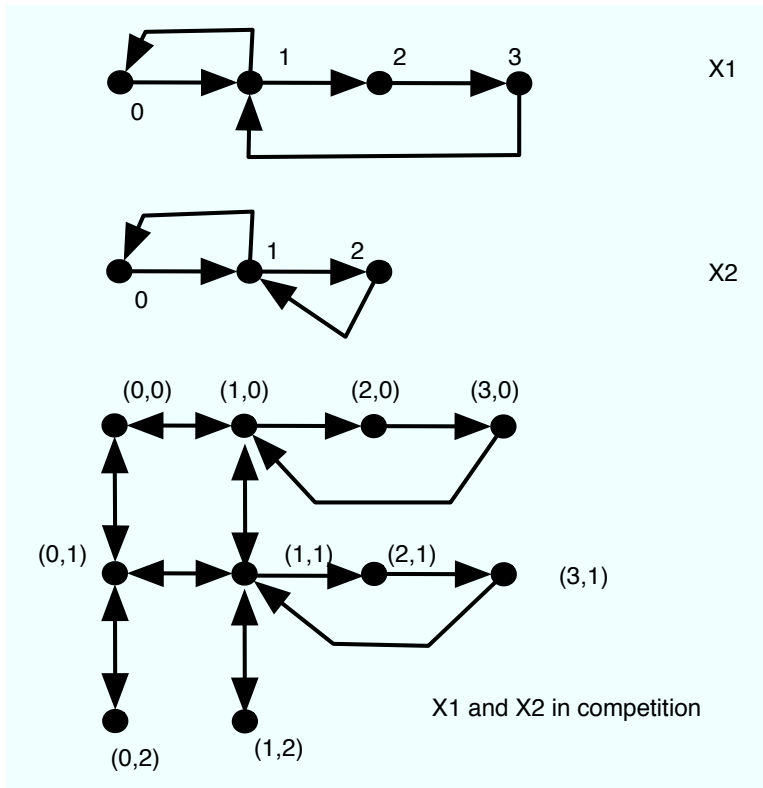
$$(v \otimes w) \times (A(\mathcal{B}) \oplus_g B(\mathcal{A})) = 0.$$

- Thus the result for CT-SAN is simpler (no restriction on the FDG)

## Competing Markov chains in Discrete Time: Simple

- Chain  $X1$  and  $X2$  compete over one resource,
- if  $X1$  has the resource, all the transitions of chain 2 are cancelled except self loops.
- if  $X1$  does not own the resource,  $X2$  evolves independently.
- $X2$  cannot block  $X1$ .
- When  $X2$  owns the resource,  $X1$  can move and if it takes the resource  $X2$  is now blocked.

# Competition in CT/DT



## Differences between CT and DT

- CT models of competition have been proved to have product form (Boucherie 94).
- Strict priority (with preemption to take the resource) in DT.
- Race in CT.
- Priority implies (Functional Dependency graph = DAG)
- Several movements in DT, only one in CT.
- Cancellation of states in CT, open question in DT.
- Product form in both cases.

## A more complex model of Competition/Collaboration

- H1': The resources are distributed among the  $X_i$  at every time slots.
- H2': If  $X_i$  receives  $b$  resources, it performs  $b$  steps of the Markov chain
- H3': Priority based allocation.  $X_1$  has the highest priority.  $X_1$  receives one resource.
- H4': According to the states of  $X_1..X_{k-1}$ ,  $X_k$  receives  $b$  resources, the transition matrix of  $X_k$  is  $(M_k)^b$
- H5': When no resources are given to  $X_k$ , it is blocked. Its transition matrix is  $Id$ .

**Theorem 4** *Consider a collection of  $N$  chains  $X_1, \dots, X_N$  in competition over a set of  $R$  equivalent resources. Suppose that assumptions  $H1'$  to  $H5'$  are satisfied. Assume that the Markov chain of the DTMC modeling the competition is ergodic, then the steady-state distribution has product form.*



## Part IV : Continuous-Time

### Domino Synchronizations, no functions

$$Q = \bigoplus_{i=1}^N Q_l^i + \sum_s \bigotimes_{i=1}^N Q_s^{(i)} + D$$

## Domino Synchronization



- An **Ordered List of Automaton**.
- Domino effect: when a domino tile falls down, the next tile does the same (success)
- But the tile may stay up and the automata in the remaining part of the list does not move.

## Domino Synchronization



Figure 5: Left: success: all the tiles are down; Right: the synchronization fails: the green tile is still up at the end.

## Some definitions

- Here we only consider a domino synchronization with 3 automata.
- The automata are called the master, the slave and the relay.
- The master initiates the synchronization.
- The slave obeys and the synchronizations propagates to the relay, or the slave refuses and the relay does not move (but the master has changed its state).
- Finally the relay obeys (or maybe not).

## Matrix Description

- A domino synchronization is described by 3 matrices.
  - Master description:  $M^{(r)}$ , a transition rate matrix.
  - Slave description:  $E^{(r)}$ , a transition probability matrix such that
    - Either  $E^{(r)}[k, k] = 0$  (the synchro succeeds at this step) or  $E^{(r)}[k, k] = 1$  (the synchro fails at this step).
    - If  $E^{(r)}[k, k] = 0$  row  $k$  of  $E^{(r)}$  gives the transition probability out of state  $k$  for the slave.
    - Relay description:  $T^r$ , a transition probability matrix (same assumptions as  $E^{(r)}$ ).
    - $E^{(r)}$  is decomposed into  $E_1^{(r)}$  (the slave obeys) and  $E_2^{(r)}$  (the slave refuses)
- Local Transitions: transition rate matrices.

## Normalization

- All matrices are normalized, i.e. for all  $k$  we have:

$$M^{(r)}[k, i] \geq 0 \text{ if } i \neq k \text{ and } \sum_i M^{(r)}[k, i] = 0,$$

$$E^{(r)}[k, i] \geq 0 \text{ and } \sum_i E^{(r)}[k, i] = 1,$$

$$T^r[k, i] \geq 0 \text{ and } \sum_i T^r[k, i] = 1.$$

- Normalization of the SAN: based on the normalization of the synchronizations ( $N^r$ ) as the local transitions are already associated to transition rate matrices.
- Definition: Let  $M$  be a matrix,  $\sigma(M)$  is a diagonal matrix with the size of  $M$  such that for all index  $i$ ,  $\sigma(M)[i, i] = \sum_j M[i, j]$ . As usual  $diag(M)$  is a diagonal matrix whose elements are the diagonal elements of  $M$ .

## A normalized tensor representation of a Domino

1.  $(M^{(r)} - \text{diag}(M^{(r)})) \otimes E_1^{(r)} \otimes T^r \otimes I_1$ : the slave accepts the synchronization.
2.  $(M^{(r)} - \text{diag}(M^{(r)})) \otimes E_2^{(r)} \otimes I \otimes I_1$ : the slave does not accept the synchronization.
3.  $\text{diag}(M^{(r)}) \otimes \sigma(E_1^{(r)}) \otimes \sigma(T^r) \otimes I_1$ : normalization of term 1.
4.  $\text{diag}(M^{(r)}) \otimes \sigma(E_2^{(r)}) \otimes I \otimes I_1$ : normalization of term 2.

## Main Result

Theorem: Consider a SAN with  $n$  automata and  $s$  with Domino synchronizations. Consider matrices  $\overline{M}^{(r)} = M^{(r)} - \text{diag}(M^{(r)})$  and  $E^{(r)}$  associated to the description of synchronization  $r$ . Let  $g_l$  an eigenvector of both  $E^{(r)}$  and  $\overline{M}^{(r)}$ . We assume that  $g_l$  exists. Let  $\Omega_r$  (resp.  $\Gamma_r$ ) be the eigenvalue for matrix  $E^{(r)}$  (resp.  $\overline{M}^{(r)}$ ) associated to  $g_l$ . If  $g_l$  is in the kernel of matrix

$$A_l = F_l + \sum_{r=1}^R \left( M^{(r)} 1_{msr(r)=l} + \Gamma_r (E^{(r)} - \sigma(E^{(r)})) 1_{sl(r)=l} + \Gamma_r \Omega_r (T^r - \sigma(T^r)) 1_{rl(r)=l} \right)$$

then the steady-state distribution has a product form solution:

$$Pr(X_1, X_2, \dots, X_n) = C \prod_{l=1}^n g_l(X_l), \quad (5)$$

and  $C$  is a normalization constant.



## Sketch of the Proof-1

Main Idea: ALGEBRA...

$(\otimes_l g_l) \times (\oplus_l F_l + \sum_i \otimes_l B_l^i)$	Reorganize the product (distributivity)
$\sum_i (\otimes_l g_l) \times (\otimes_l M_l^i)$	Then use compatibility with $\times$
$\sum_i (\otimes_l (g_l \times M_l^i))$	Assumptions on eigenvectors
	Factorize and find some matrices $P_l^i$
$\sum_i (\otimes_l (g_l \times P_l^i))$	such that
$\pi_{j_0} P_{j_0}^i = 0$	for all $i$ , it exists $j_0$ (automaton index)

## Sketch of the Proof-2

The description of  $(\pi_1 \otimes \pi_2 \otimes \dots \otimes \pi_n)Q$  consists in 7 terms (three coming from the tensor sum, two for the Domino and two for the normalization of the Domino):

$$\begin{aligned} & (\pi_1 F_1 \otimes \pi_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 F_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 \otimes \pi_3 F_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 (M^{(r)} - \text{diag}(M^{(r)})) \otimes \pi_2 E_1^{(r)} \otimes \pi_3 T^r \otimes \dots \otimes \pi_n) \\ + & (\pi_1 (M^{(r)} - \text{diag}(M^{(r)})) \otimes \pi_2 E_2^{(r)} \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \text{diag}(M^{(r)}) \otimes \pi_2 \sigma(E_1^{(r)}) \otimes \pi_3 \sigma(T^r) \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \text{diag}(M^{(r)}) \otimes \pi_2 \sigma(E_2^{(r)}) \otimes \pi_3 \otimes \dots \otimes \pi_n) \end{aligned}$$

## Sketch of the Proof-3

- $\pi_1(M^{(r)} - \text{diag}(M^{(r)})) = \pi_1\Gamma_r.$
- $\sigma(T^r) = I$
- $\sigma(E_1^{(r)}) + \sigma(E_2^{(r)}) = I$
- After simplification:

$$\begin{aligned} & (\pi_1 F_1 \otimes \pi_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 F_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 \otimes \pi_3 F_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \Gamma_r \otimes \pi_2 E_1^{(r)} \otimes \pi_3 T^r \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \Gamma_r \otimes \pi_2 E_2^{(r)} \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \text{diag}(M^{(r)}) \otimes \pi_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \end{aligned}$$

## Sketch of the Proof-4

- Many algebraic manipulations (in the paper)...
- The ordinary product is compatible with the tensor product (i.e.  $(\lambda A) \otimes B = A \otimes (\lambda B)$ ).
- $\sigma(E_2^{(r)}) = E_2^{(r)}$  and  $\sigma(E_1^{(r)}) + \sigma(E_2^{(r)}) = I$ . Using the distributivity, after cancellation we get:

$$\begin{aligned} & (\pi_1(F_1 + M^{(r)}) \otimes \pi_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2(F_2 - \Gamma_r \sigma(E_1^{(r)})) \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 \otimes \pi_3 F_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \Gamma_r \otimes \pi_2 E_1^{(r)} \otimes \pi_3 T^r \otimes \dots \otimes \pi_n) \end{aligned}$$

## Proof

- We apply the assumption on the eigenvalue of  $E_1^{(r)}$ .
- After substitution:

$$\begin{aligned} & (\pi_1(F_1 + M^{(r)}) \otimes \pi_2 \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2(F_2 + \Gamma_r(E^{(r)} - \sigma(Er))) \otimes \pi_3 \otimes \dots \otimes \pi_n) \\ + & (\pi_1 \otimes \pi_2 \otimes \pi_3(F_3 + \Gamma_r\Omega_r(T^r - \sigma(T^r))) \otimes \dots \otimes \pi_n) \end{aligned}$$

- Now proceed with the other synchronizations...
- We obtain matrix  $A_l$
- As  $\pi_l A_l = 0$ , we get  $(\pi_1 \otimes \dots \otimes \pi_l \otimes \dots \otimes \pi_n)(Id \otimes A_l \otimes Id) = 0$ .
- And the proof is complete.

## Example-1: G-network with trigger

- Infinite state space.
- Each automaton models the number of positive customers in a queue.
- The local transitions are the external arrivals (rate  $\lambda_l$ ) and the departures to the outside (rate  $\mu_l$  multiplied by probability  $d_l$ ).
- Synchronization : departure of a customer on the master (the end of service with rate  $\mu_l$  and probability  $(1 - d_l)$ ), the departure of a customer on the slave (a customer movement, if there is any), the arrival of a customer on the relay (always accepted).

## Example-1-Matrix

$$\begin{aligned}Q^{(l)} &= \lambda_l(U - I) + \mu_l d_l(L - I^0), \\M^{(r)} &= \mu_l(1 - d_l)(L - I^0), \\E^{(r)} &= L \text{ and } T^r = U.\end{aligned}\tag{6}$$

- Tridiagonal matrix.
- $\pi_l$  has a geometric distribution with rate  $\rho_l$ :

$$\rho_l = \frac{\lambda_l + \sum_{r=1}^R \Omega_r \Gamma_r 1_{rl(r)=l}}{\mu_l + \sum_{r=1}^R \Gamma_r 1_{sl(r)=l}}.$$

- Because of its geometric distribution,  $\pi_l$  is an eigenvector of operators  $\overline{M^{(r)}}$  and  $E^{(r)}$ .
- Finally:  $\Omega_r = \rho_{sl(r)}$  and  $\Gamma_r = \rho_{msr(r)} \mu_{msr(r)} (1 - d_{msr(r)})$ ,

## Example 2-Three Deletions

- The deletion in the relay only occurs if the deletion in the slave was successful.

$$\begin{aligned}
 Q^{(l)} &= \lambda_l(U - I) + \mu_l d_l(L - I^0), \\
 M^{(r)} &= \mu_l(1 - d_l)(L - I^0), \\
 E^{(r)} &= T^r = L.
 \end{aligned} \tag{7}$$

- Tridiagonal matrix.
- $\pi_l$  has a geometric distribution with rate  $\rho_l$ :

$$\rho_l = \frac{\lambda_l}{\mu_l + \sum_{r=1}^R \Omega_r \Gamma_r 1_{rl(r)=l} + \sum_{r=1}^R \Gamma_r 1_{sl(r)=l}}.$$

and  $\Omega_r = \rho_{sl(r)}$  and  $\Gamma_r = \rho_{msr(r)} \mu_{msr(r)} (1 - d_{msr(r)})$ .



## Conclusion

- New results for Product Form for steady state distribution of multicomponent CTMC and DTMC
- As usual, DT is harder than CT
- Tensor based proof.
- Generalization to SAN with functions and synchronizations.
- Many competition/collaboration rules with product form.
- Links between partial balance and definition of  $\mathcal{S}(\alpha)$ .
- Generalization to SAN with functions and synchronizations.